

# *Implementation of Quality-of-Security-Service in communication structure for 3D-MPSoCS Protection*

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# SUMMARY

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1. Introduction
2. 3D-MPSoCs
3. HoCs: 3D Communication structure
4. Our Work

Goal 1: Mechanisms to support QoSS.

Goal 2: Evaluation of HoC performance.

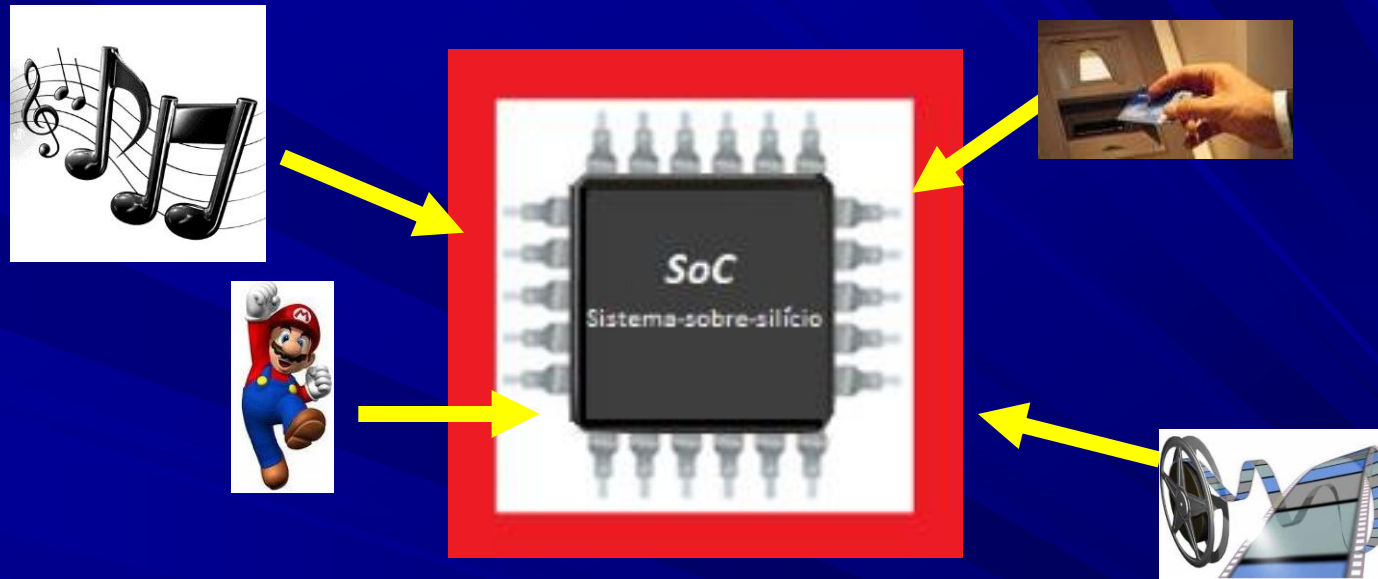
5. Results
6. Conclusions

# MOTIVATION



- To integrate more functionality into smaller devices.
- To increase performance, reduce costs.

# MOTIVATION

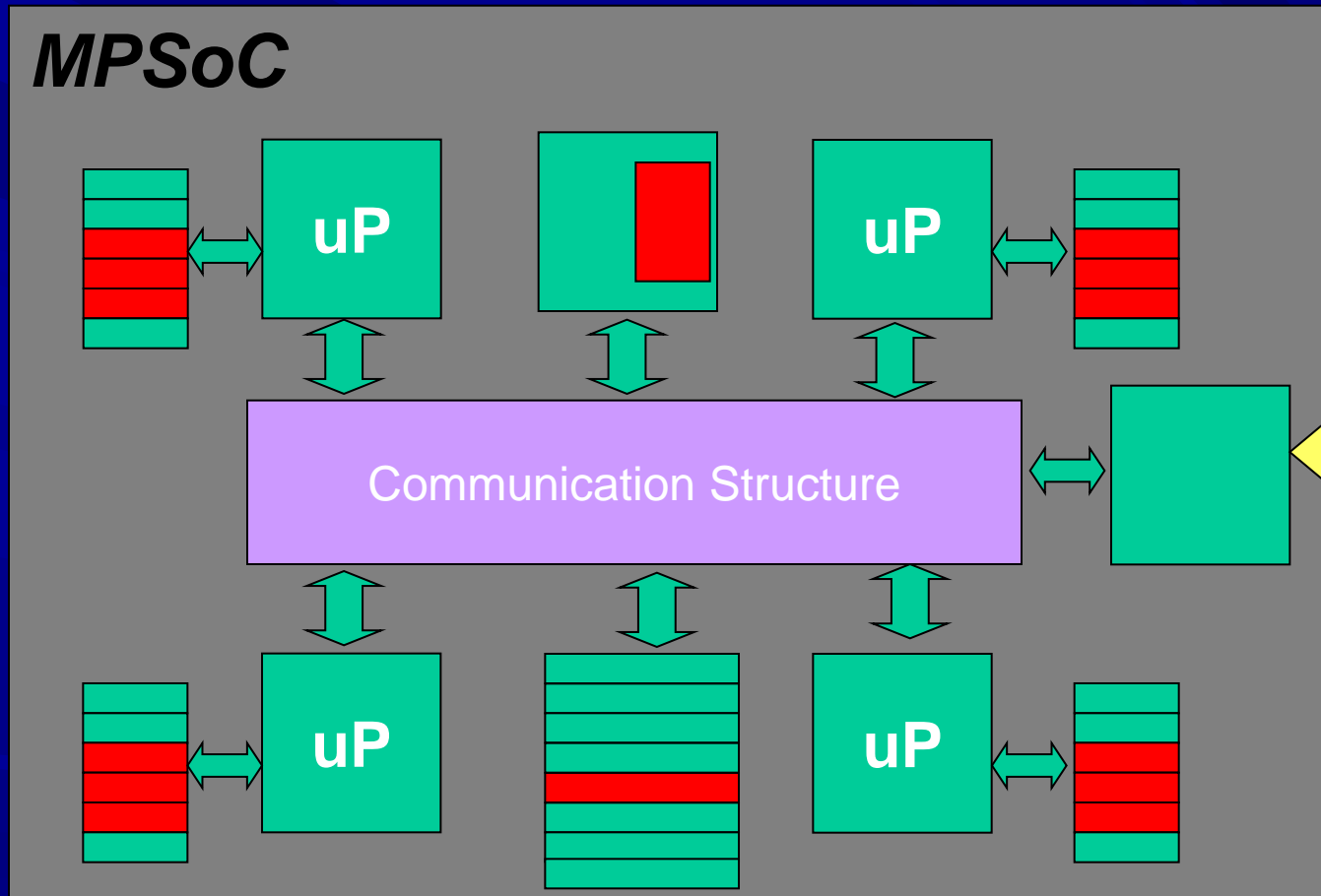


- Cost effective:
- \* General purpose SoC (*MPSoCs*).
  - \* Integrate different applications on the same chip.

Applications: Communication requirements and design constraints.

## MULTI-APPLICATION SYSTEM

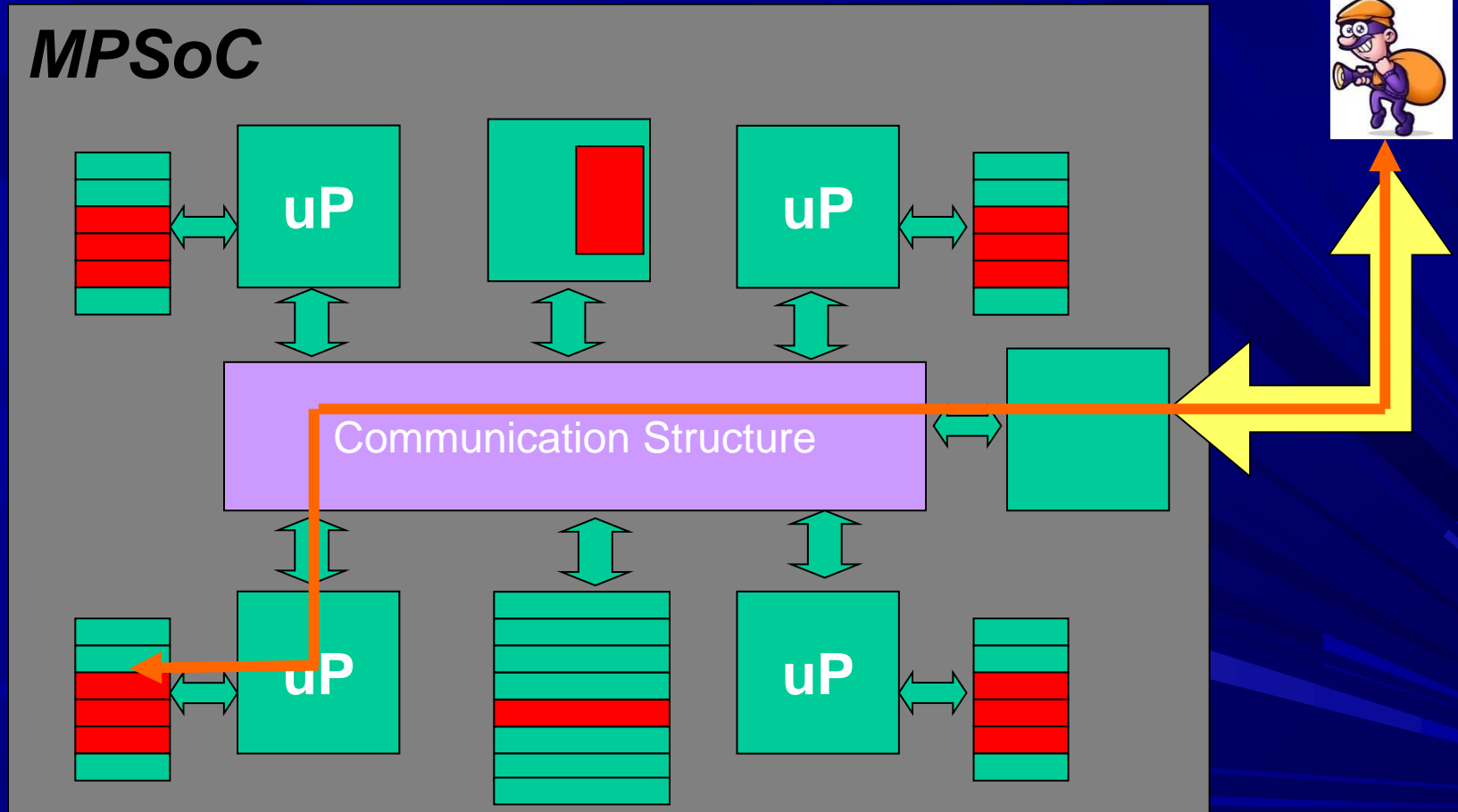
# Problem



Software attacks!

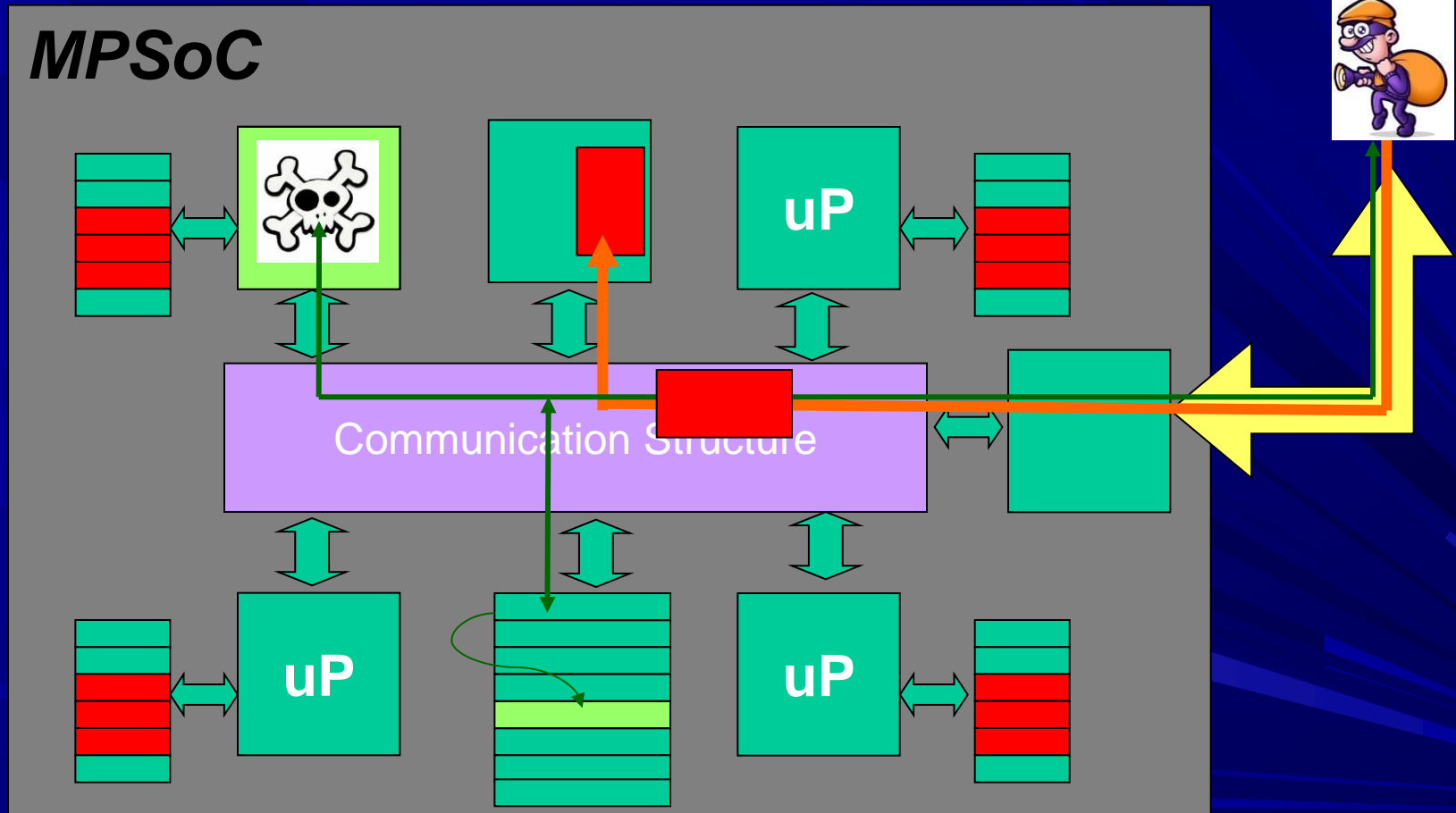
- Security incidents: 80% via **software**.

# Problem



Explore the SoC vulnerabilities.

# Problem



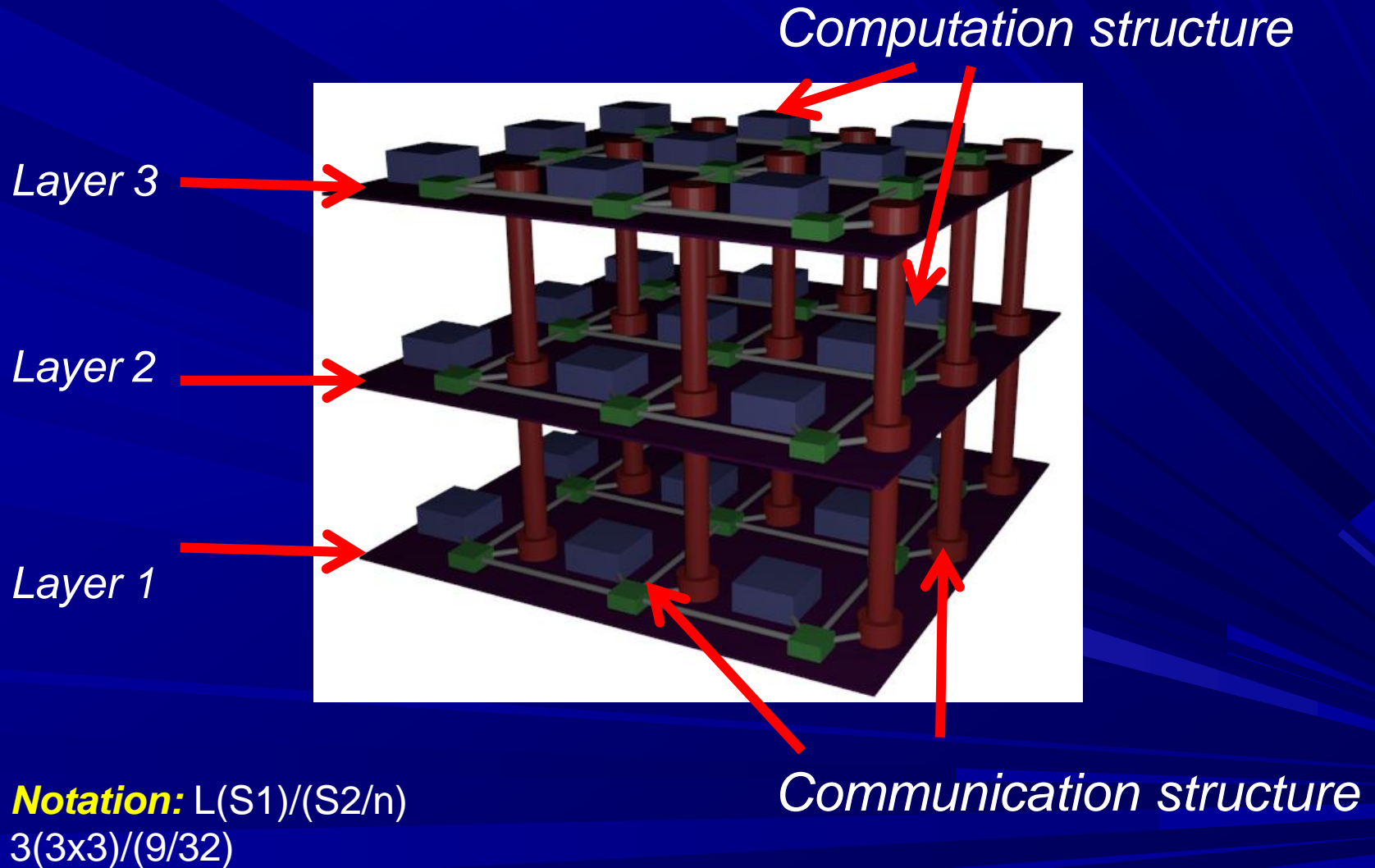
Infection: Takes advantage of the trusty component's rights!!

# 3D-MPSoCs

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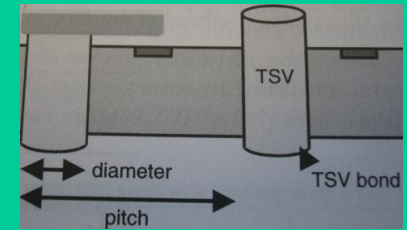


# 3D-MPSoC

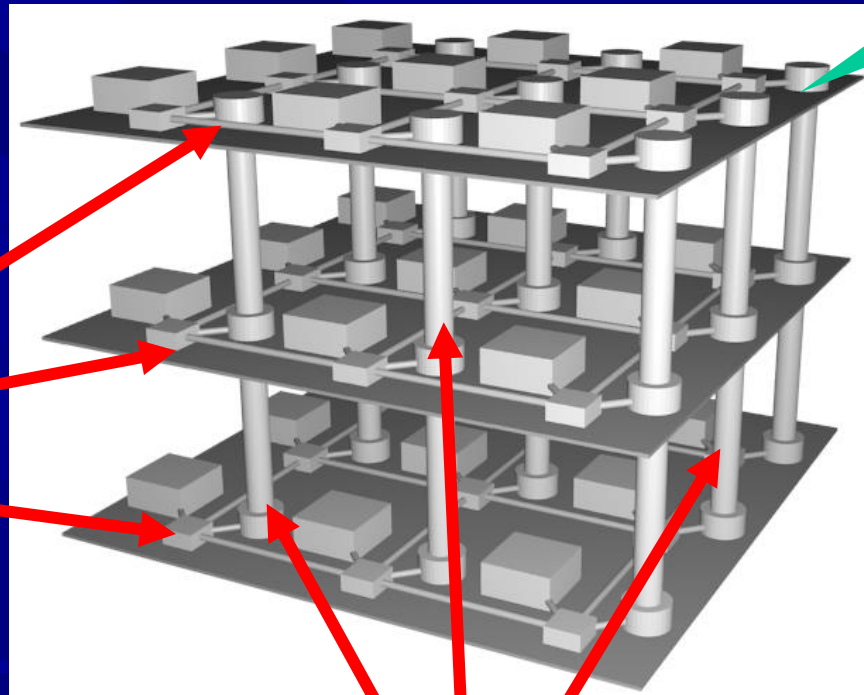


# HoCs: 3D-MPSoC Communication Structure

**HoCs:** Hybrid-On-Chip CS



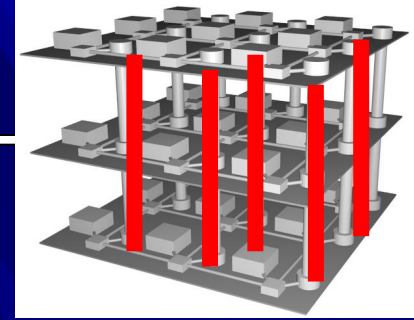
NoCs



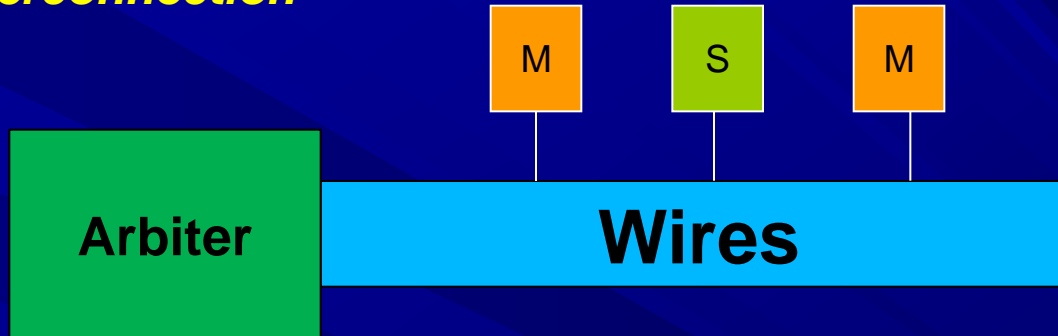
Bus

- Short connections.
- Low capacity.
- High frequency.
- Defects.
- Area consumers.

# HoCs: Bus

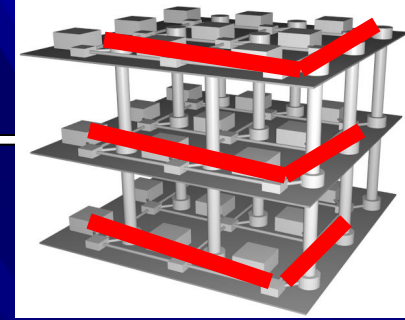


## *Vertical interconnection*

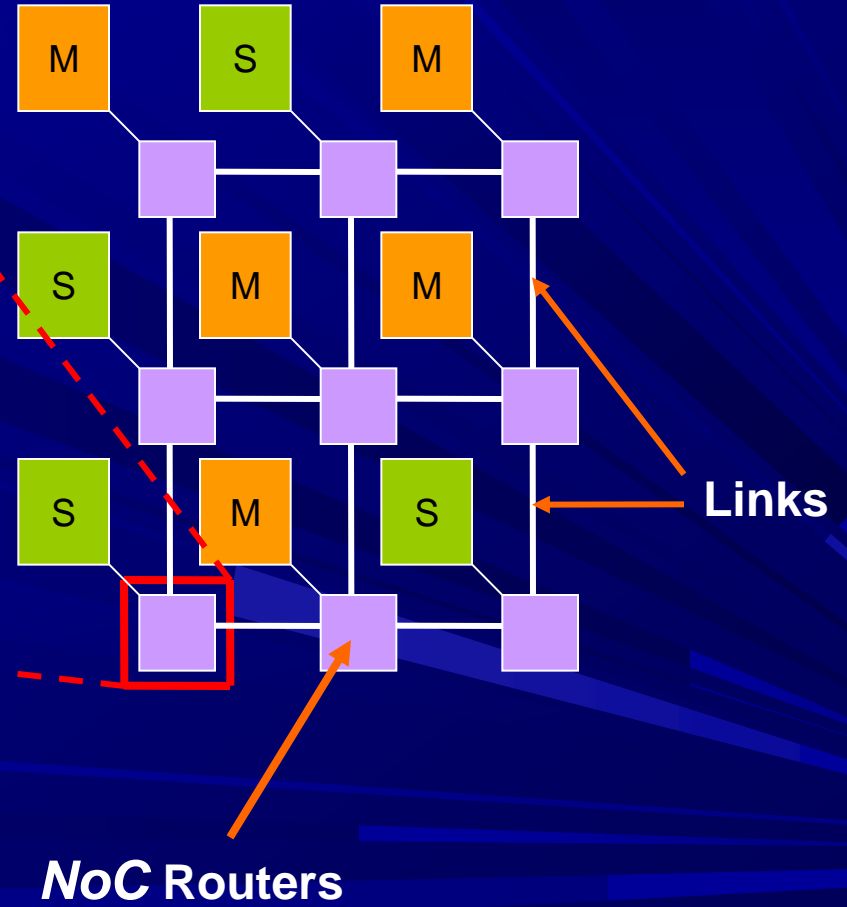
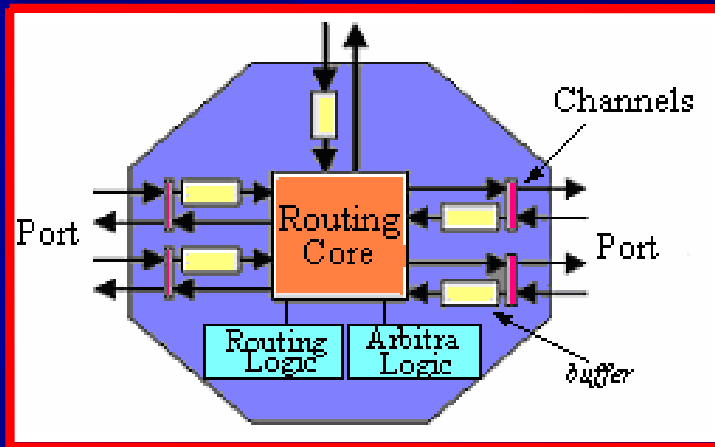


- Low cost CS with predictable latency.
- Not scalable.
- Number of interlayer links (performance/cost-reliability)
  - Higher: Improve performance of the system.
  - Lower: Prone to defects.

# HoCs: NoCs

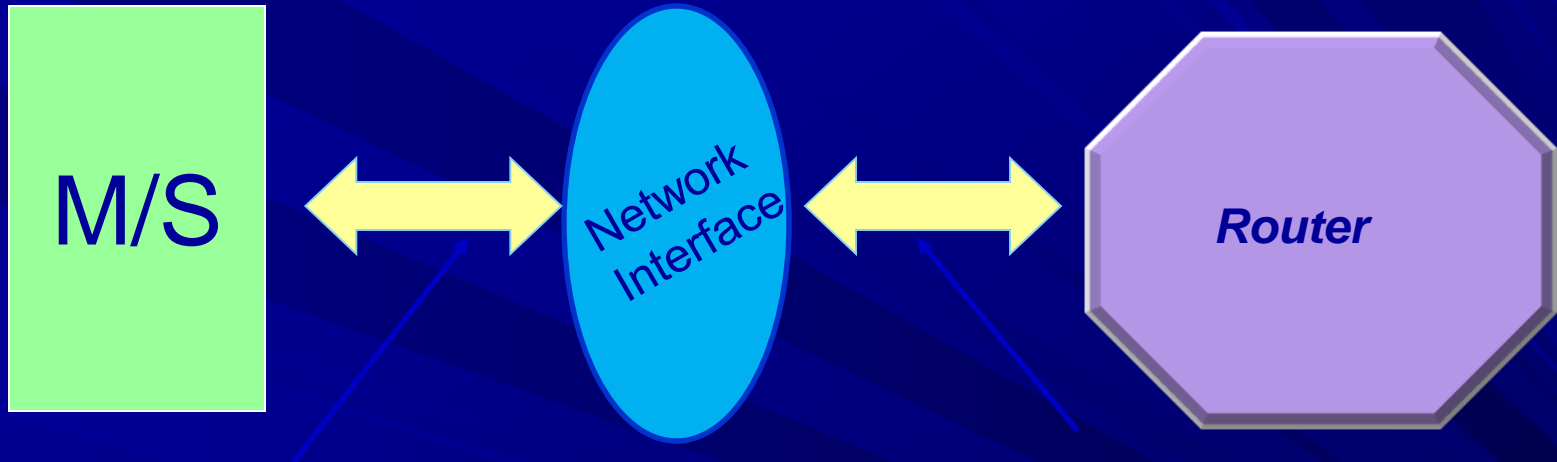
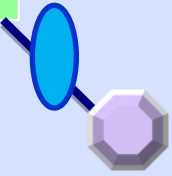


## Horizontal interconnection



# HoCs: NoCs

M/S



*Network Protocol*

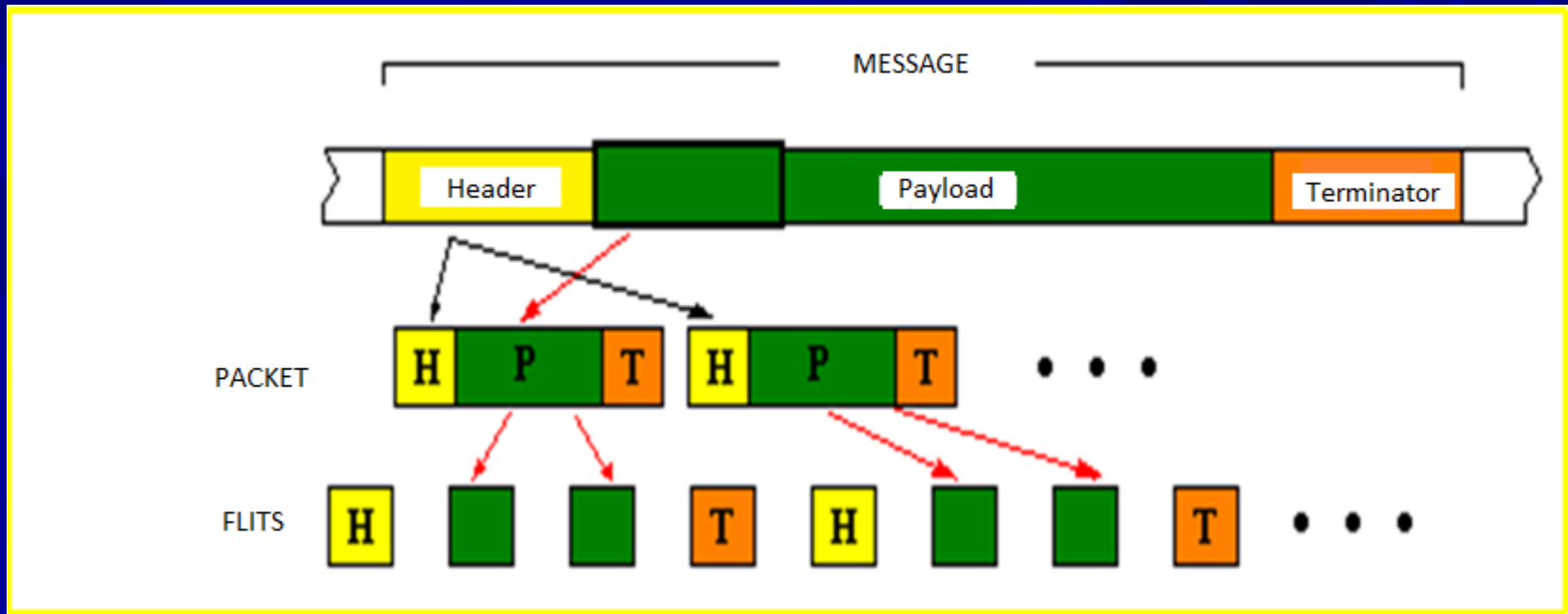
## Source

- \* Accesses routing tables.
- \* Assembles packets.
- \* Splits into flits.

## Destination

- \* Synchronizes.
- \* Drops routing information.

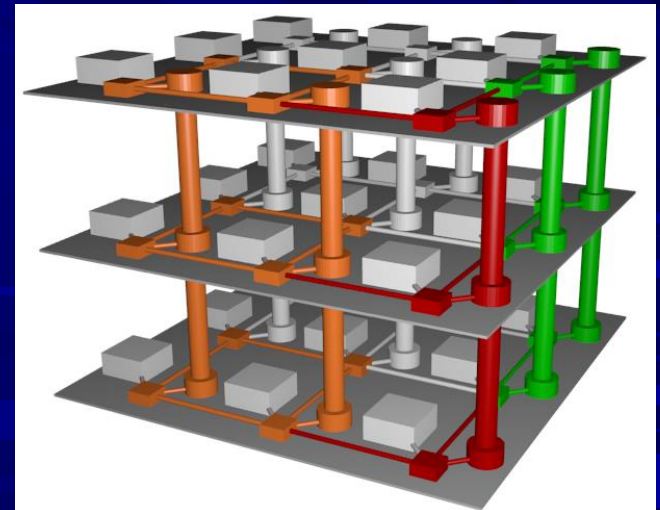
# Communication



Quality (QoS)

Security (S)

QoSS (QoS)

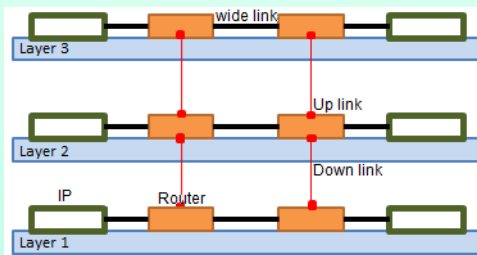
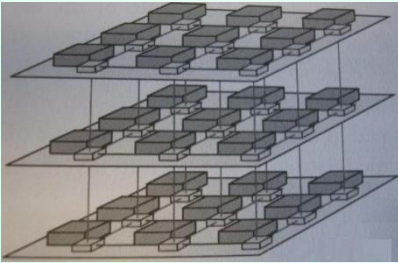


# Challenges

# 1. Efficiency

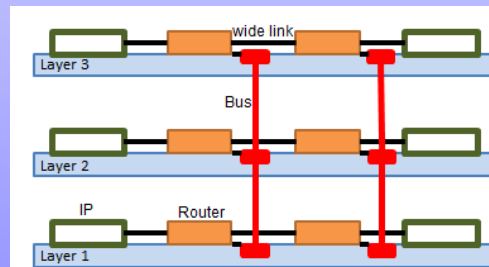
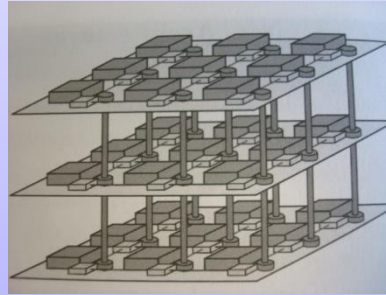
- CS is the bottleneck of the 3D-MPSoC.
- Several works address the design of 3D-CS.

## Single 3D



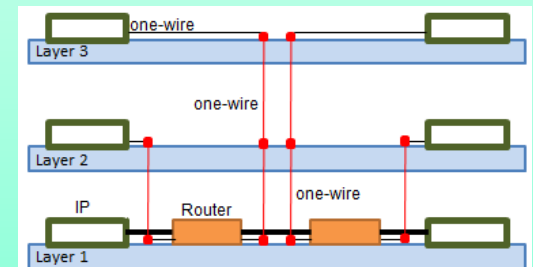
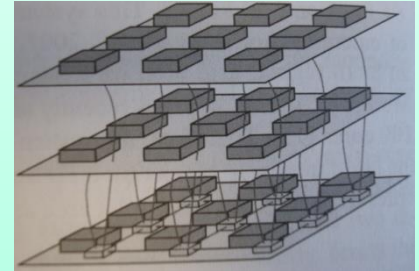
[HEA10]

## Stacked



[SHE10]

## Ciliated

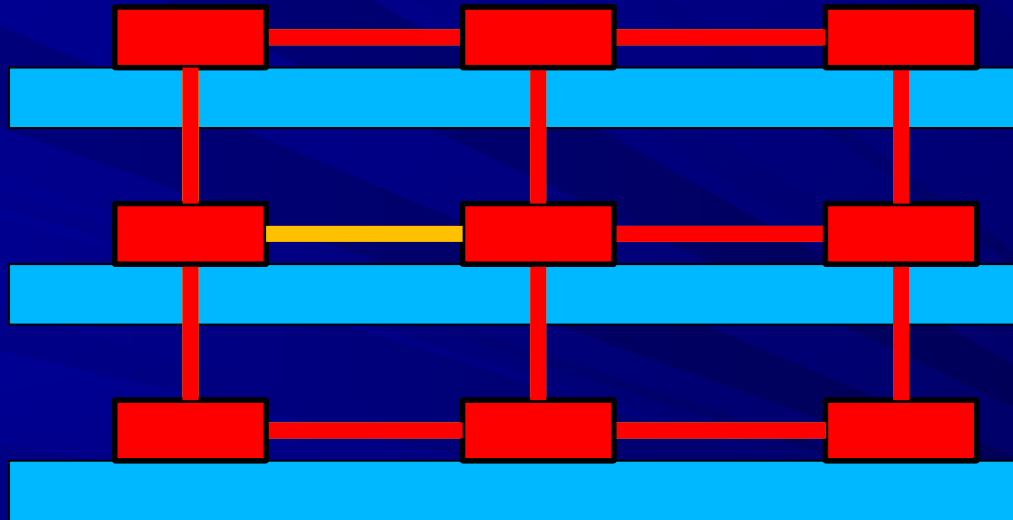


[STA11]

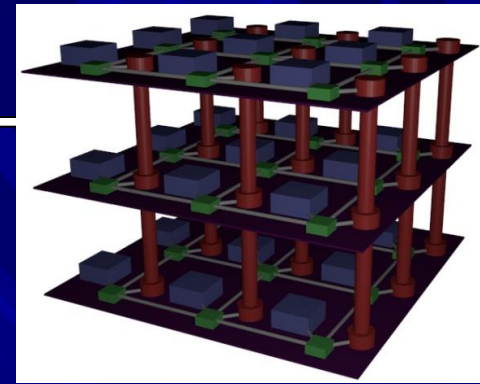
**BEST EFFORT ARCHITECTURES!  
WITHOUT SECURITY**



## 2. Security



# 3D-MPSoC characteristics



- **Multi-application**
  - Different
    - **Functional/Communication requirements.**
    - **Security requirements** (multi security-policies).
- **Dynamicity**
  - Applications may change (**dynamic security requirements**).
  - New applications may have
    - Tighter communication requirements.
    - Stronger/weaker security requirements.
- **Heterogeneity**
  - Components with different performance.
  - From different providers (are they trustworthy?).
- **Observability**
  - Track of critical information (i.e. state of IPs for tasks migration).

# Dynamic security requirements

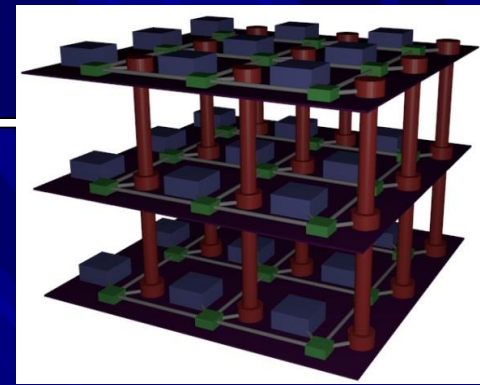
- The security policy of the 3D-SoC can change as a consequence of three factors:
  - ***New application*** is mapped on the 3D-SoC.
  - ***Current application is reallocated*** on the 3D-SoC (i.e. Task migration).
  - ***New 3D-SoC operation scenario***.

***Islands:*** IPs or clusters of IPs.

# 3D-HoC services

- Just an extension of 2D?

**NO**



- **3D presents new challenges**
  - All get worst: multi-application, dynamicity, heterogeneity.
  - Increase of faults (TSVs and thermal effects).
- **3D presents new opportunities:**
  - Promote design strategies (prohibitive in performance at 2D-SoC)
    - Huge amount of task migration.
    - Layers specialization.
    - Cluster-style design (clusters linked through a 3D-HoC).
  - Huge set of configuration parameters
    - Computation structure
    - Communication structure

# Opportunities

# Security Opportunities

- **COMPUTATION STRUCTURE:**

- High level of integration: More IPs integrated to the 3D-MPSoC can be dedicated to security.
  - Cryptoprocessors
  - Security IPs.

- **COMMUNICATION STRUCTURE:**

- 3D-MPSoCs are foreseen as communication-centric systems.
- All software attacks start with an abnormal communication.
- Main role of the CS in the system operation can be used for detect an attack.

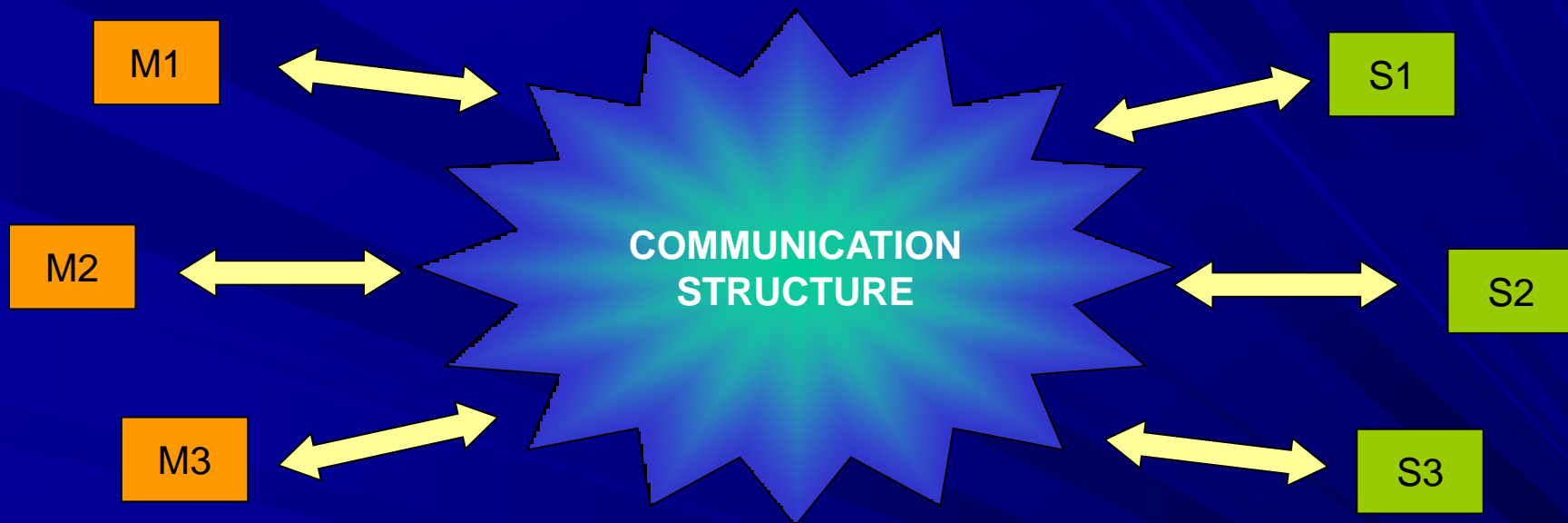
# OUR WORK

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## *Goal:*

1. To integrate security mechanisms to the HoC in order to provide different levels of security (3D-QoCS), evaluate its efficiency and efficacy.

# Communication structure



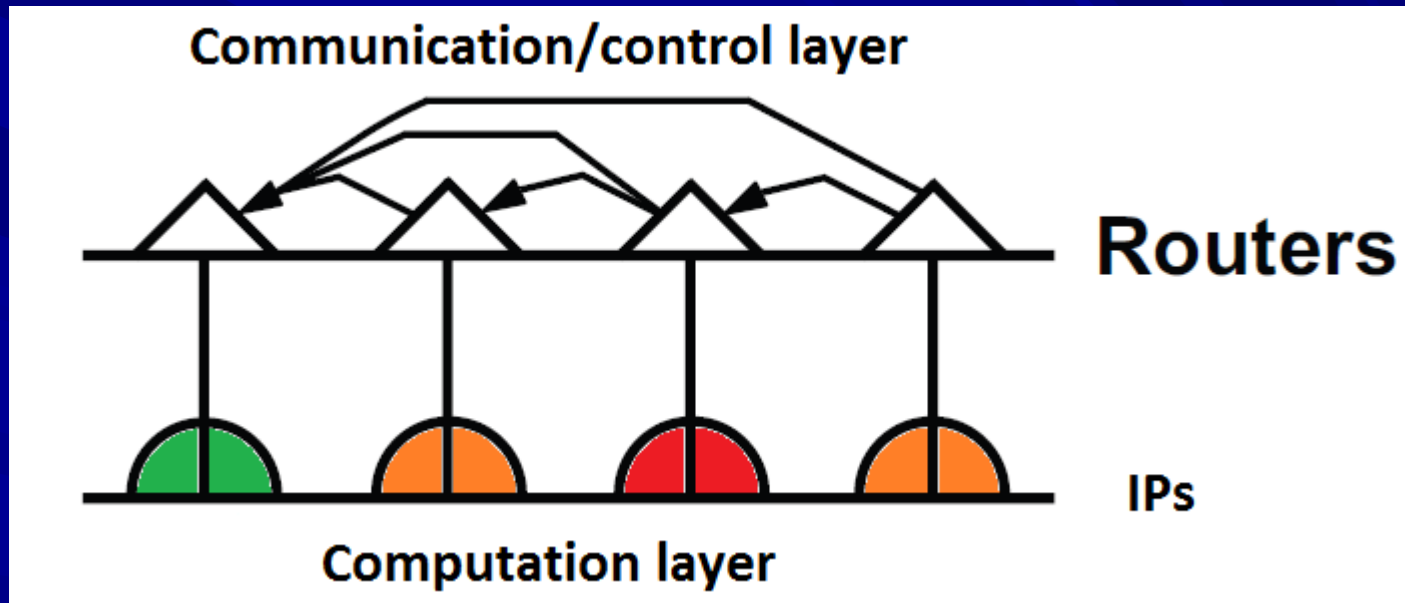
All software attack begins with an **abnormal communication**.

- Monitor information exchange.
- Detect attacks.
- Diagnosis  Trigger recovery mechanisms.



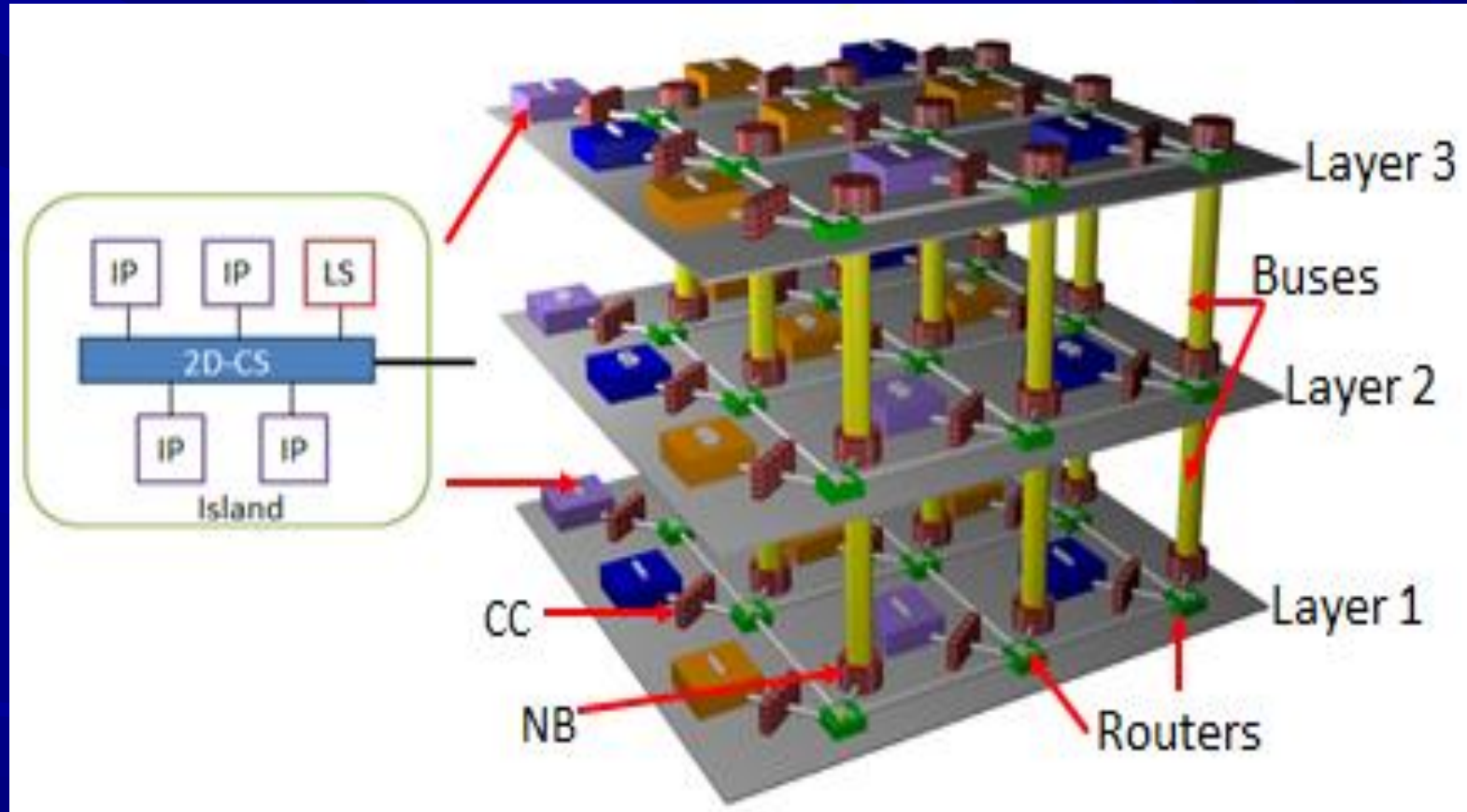
# Security Implementation

# 1. Application specific security layer



- Application specific security functionality
- Isolation
- Passive monitoring
- Layers can be fabricated at different foundries and integrated in a third trusty foundry.

## 2. Split security at all the layers



*Islands:* IPs or clusters of IPs.

# 2. Split security at all the layers

## *Characteristics:*

We implement two security services at the 3D-HoC:

- i) *authentication*: verifying the source integrity.
  - ii) *access control*: certifying the authorized use of the system.
- Different security choices (L0- L3 ):
    - Special configuration of the *security mechanism*.
    - Higher security may imply in higher costs.
    - Selection of a security level:
      - Security requirements of the system.
      - Resources availability and cost.

3D-SoC designer may select a lower protection level in order to fulfill the performance requirements (trade-off).

# Access Control

- Place of implementation: Interface, router.
- Security levels.
- Control information: Source, type, role.

## FILTER:

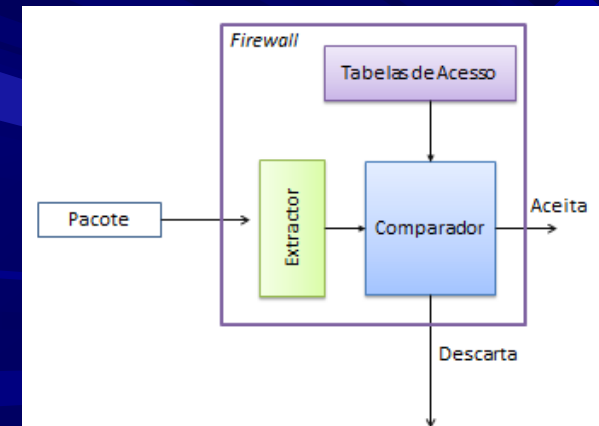
- *HoC firewall* : Allows or blocks a transaction.
- According to security policy.

## Interface:

- \* Before packet injection to the CS.
- \* Packet reception.

Access control			
	SV	TV	PV
Level 0			
Level 1	X		
Level 2	X	X	
Level 3	X	X	X

SV: Source verification.  
TV: Type verification.  
RV: Role verification.



# Authentication

- Implementation place: Interface, router.
- No cryptographic mechanisms.
- Levels of security.

## ANALIZER:

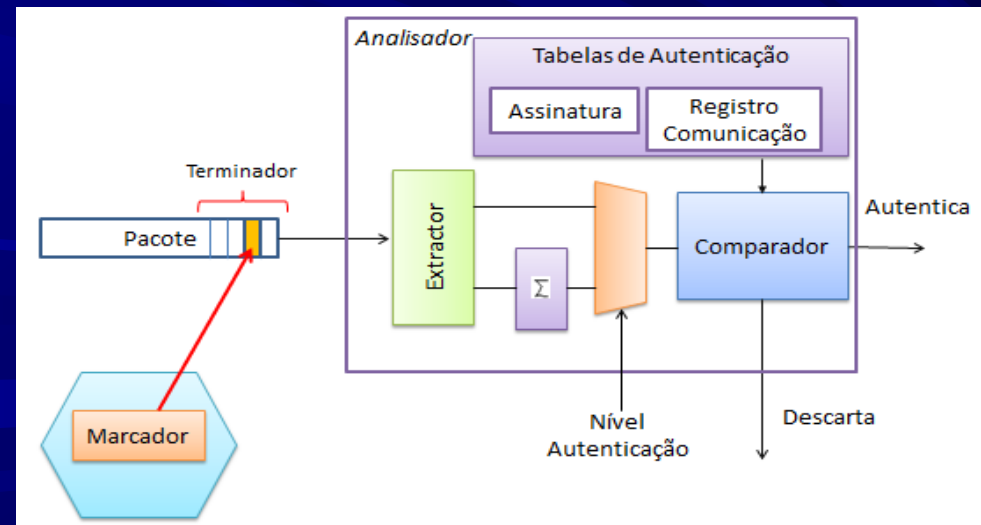
*Number of routers through the communication path.*

*Routers ID.*

*Communication code.*

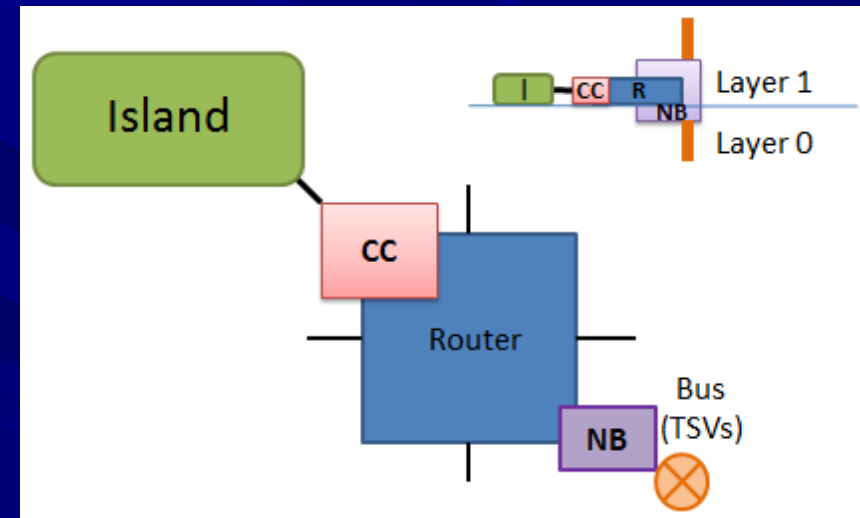
Authentication			
	NR	RP	CC
Level 0			
Level 1	X		
Level 2	X	X	
Level 3	X	X	X

NR: Router number.  
RP: Set Routers ID.  
CC: Communication code.



# 2. Split security at all the layers

- Firewalls in the 3D-HoC interfaces: Allow or block a transaction according to the matching or mismatch between the content of the packet and the security policy.
- Firewalls store the security policy information in a *security table*.
- 3D-HoCs integrates two types of interfaces:
  - Computation-Communication (**CC**).
  - NoC-Bus (**NB**).



**CC:** rules the intra-layer communication (same layer).

**NB:** rules the inter-layer communication (different layers).

SECURITY MECHANISMS

Service	Mechanism	CC	NB	L0	L1	L2	L3
AC	Destination	Island	Memory	x	x		x
	Operation	read, read-linked, write and broadcast	read, read-exclusive, write.		x	x	x
	Size	No checking	Checking		x	x	x
	Deadline/role	cycles/root-user	cycles/root-user			x	x
AU	Source	Island	Island	x	x		x
	Path	No checking	Checking		x	x	x
	ID Code	Checking	Checking			x	x

# 2. Architecture

## ***Policy keeper:***

- It stores the information of the 3D-SoC task mapping and the security policy.
- The security policy set the protection level (from L0 to L3) of each service.
- The size of the table stored by the policy keeper component depends on the number of applications, tasks and IPs integrated at the 3D-MPSoC.

## ***Reconfiguration manager:***

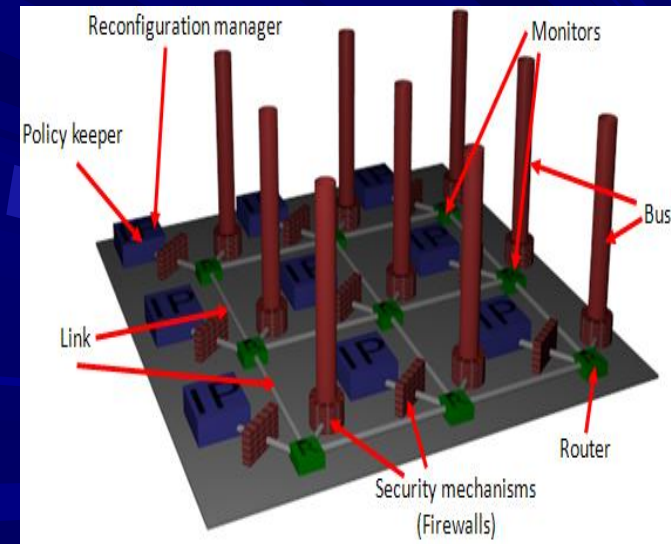
- Coordinates the upgrading of the security table of all the firewalls.

## ***Security mechanisms:***

- Defends the 3D-MPSoC against possible attacks.
- Uses the information embodied in the packets.
- Able to be upgraded.

## ***Monitor:***

- Audits the communication behavior of the 3D-SoC.
- Determine the completion of the transaction.
- Embodied at the routers of the 3D-HoC.





# 2. Functionality

## 1. Analysis the security policy

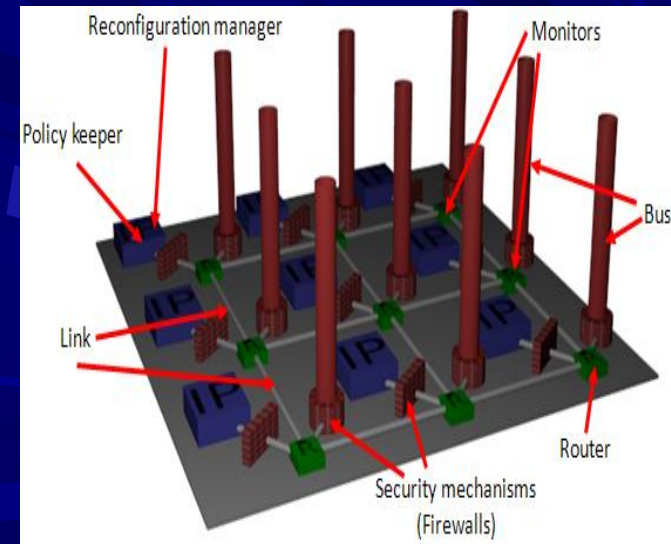
- Identify the firewalls that must be configured (target firewall).
- Which, where, new data.

## 2. Configuration of security mechanisms

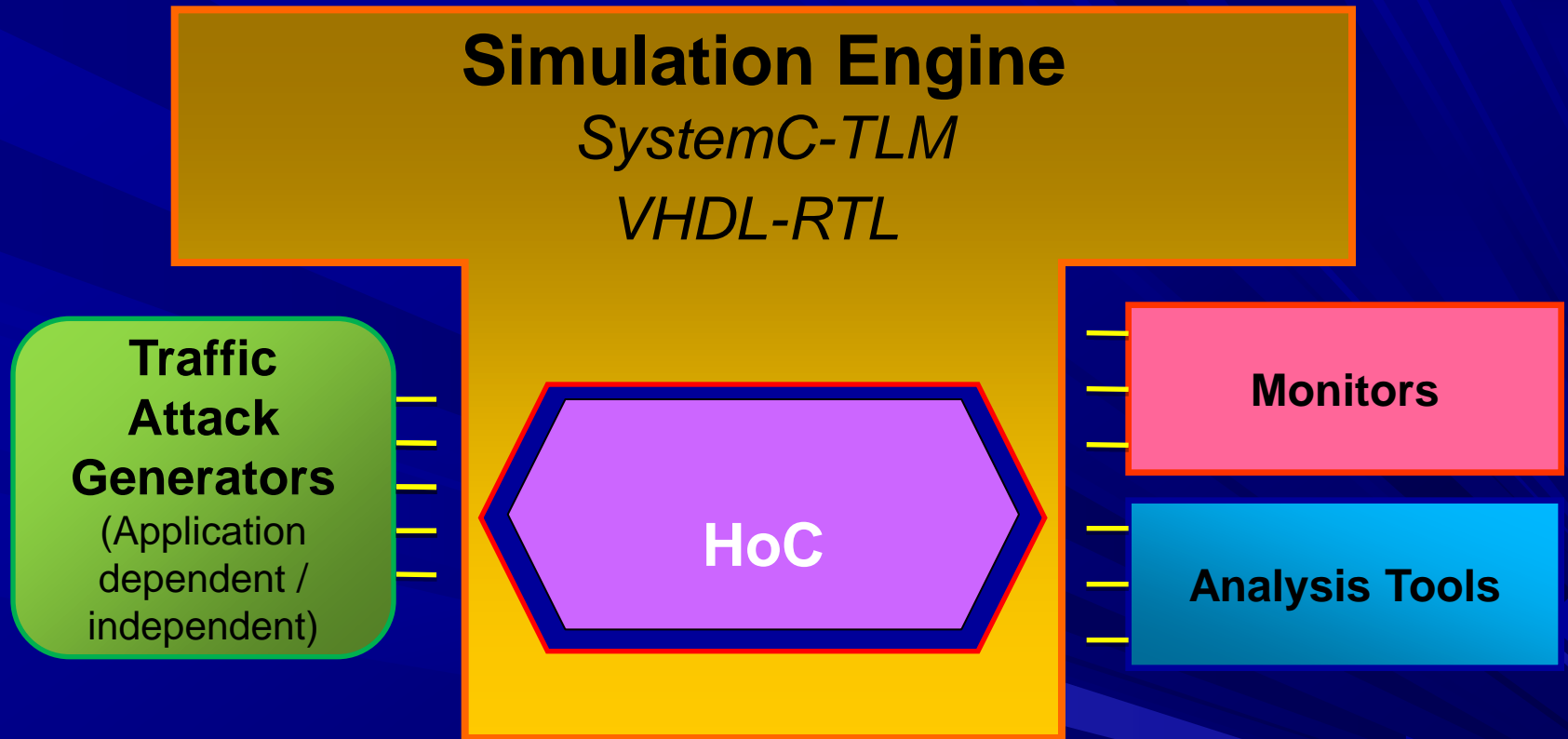
- Block injection of new data whose destination is linked to the target firewall.
- Send new data (local and global configuration).

## 3. Recovery

- Unblock communication.
- Resume operation.



# Evaluation



HoC simulation and evaluation framework.  
Supports different traffic conditions.

# Experimental Setup

**CS:** *2D-NoC (application specific layer)*

*HoCs (security in all the layers)*

## **HoC Configuration**

- Stacked, single, ciliated and **3D-HoC**  $3(5 \times 5)/25/32$
- XYZ routing algorithm
- 75 IP cores 3D-MPSoC
- Round-Robin
- Simple/QoS arbiter
- FIFO memory organization

## **Simulation Conditions**

- *5 flits Payload.*
- *900.000 simulated cycles.*

# Experimental Setup

- 3 characteristics of the traffic: Nature, topology and type.
- **Topology**
  - Hot-spot
  - Transpose
  - Uniform
- Real application (3 Applications, different security policies)
- **Nature**
  - Poisson + % LRD .
- **Type of traffic**
  - Best effort
  - Priority (L M H)
  - Guarantee
- **Dynamicity** (0, 20, 40, 50, 60 ,80)

# Results

## Efficiency:

\* 3 different kind of attacks (Modification, extraction, DoS).

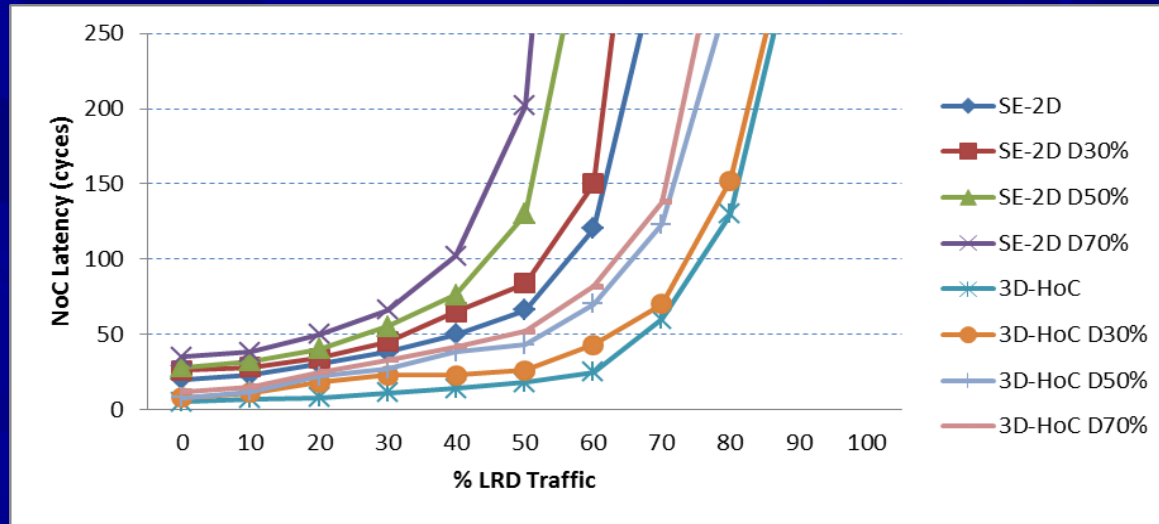
SECURITY EFFICACY		
<i>Attack scenario</i>	<i>2D-NoC</i>	<i>3D-HoC</i>
Write critical data	97%	97%
Read critical data	100%	100%
Malicious task migration	100%	100%
Nonexisting target /Repeated data	89%	89%
Communication target = source	100%	100%

- They show **identical** security efficacy (percentage of detected attacks).
- It was **expected** because the values of the security values at both alternatives were the same.
- The difference is the implementation (centralized, spread).
- 97% of efficacy mean that the **security designer should increase the protection level** in order to achieve a 100% of protection.

# Results

## Efficacy:

Latency results for CS L3 AC and AU security level and different dynamicity.

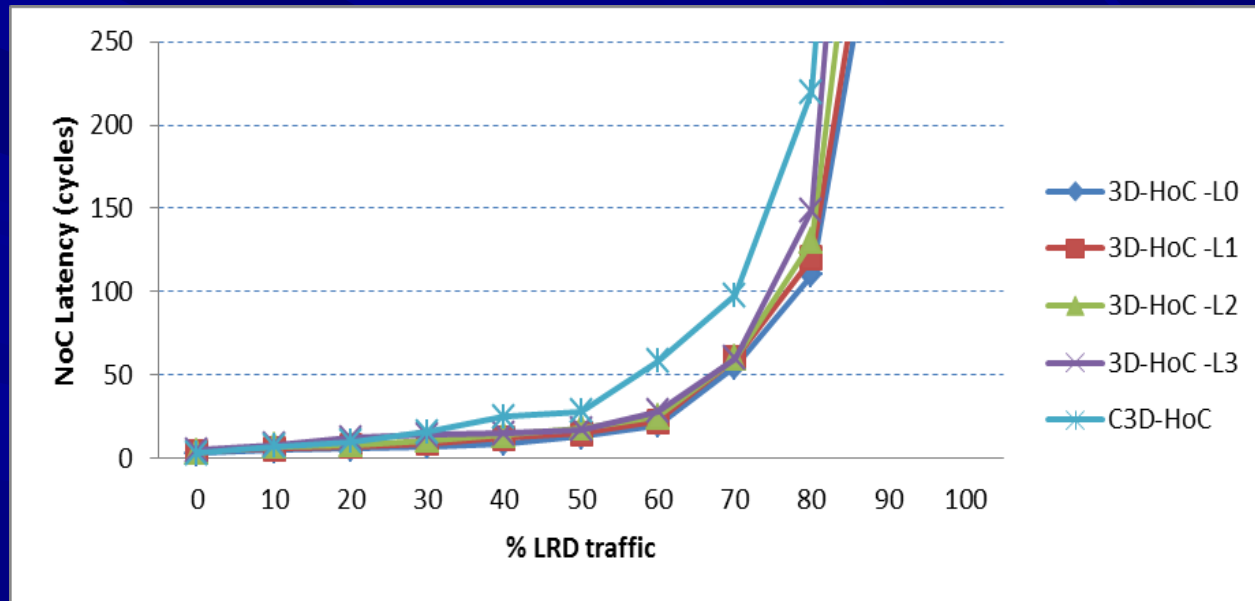


- 3D-HoC achieves a better performance when compared to 2DNoC.
- 3D-HoC is less sensible to the dynamicity of the system.
  - i) 3D technology characteristics (smaller initiator/destination paths).
  - ii) At the reconfiguration phase, only some small areas of 3D-HoC where blocked.

# Results

## Efficacy:

3D-HoC latency results for different levels of protection.



There is a **trade-off** security/performance to be explored!

# Conclusions and future work

- We propose a dynamic security enhanced 3D-HoC for 3D-SoC protection.
- We show that 3D-HoC can be an efficient structure to guarantee the protection in the system.
- 3D technology not only presents new challenges, but new opportunities to achieve a secure and efficient system.
- Three techniques are employed in order to achieve an efficient configuration:
  - Only some firewalls are upgraded, so the communication in the remaining of the system is not interrupted
  - Security customization
  - Intrinsic low latency of 3D technology.



# Conclusions and future work

- We compare our distributed architecture with a centralized one. As dynamicity increases, the distributed alternative becomes more efficient.
- As future work we plan to implement integrity and confidentiality security services.